Hash-based signatures III Stateful signatures

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SAC – Post-quantum cryptography

Merkle's (e.g.) 8-time signature system

Hash 8 one-time public keys into a single Merkle public key P_{15} .



 $S_i \rightarrow P_i$ can be Lamport or Winternitz one-time signature system. Each such pair (S_i, P_i) may be used only once.

Signature in 8-time Merkle hash tree

Signature of first message: $(sign(m, S_1), P_1, P_2, P_{10}, P_{14})$.



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Verify signature sign (m, S_1) with public key P_1 (provided in signature). Link P_1 against public key P_{15} by computing $P'_9 = H(P_1, P_2)$, $P'_{13} = H(P'_9, P_{10})$, and comparing $H(P'_{13}, P_{14})$ with P_{15} . Reject if $H(P'_{13}, P_{14}) \neq P_{15}$ of if the signature verification failed.

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Improvements to Merkle's scheme

- Each key is good only for fixed number of messages, typically 2^n .
- The public key is very short: just one hash output. But each signature contains n public keys along with the one-time signature.
- Computing the public key requires computing and storing 2ⁿ one-time signature keys.

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- Can build trees of trees where each leaf of the top tree signs the root of a tree below it. Only the top tree is needed in key generation. This increases the signature length (one one-time signature per tree) and signing time. See PhD thesis of Andreas Hülsing for an optimized schedule of what to store and when to precompute.

Stateful hash-based signatures

- Only one prerequisite: a good hash function, e.g. SHA3-512. Hash functions map long strings to fixed-length strings. Signature schemes use hash functions in handling plaintext.
- Old idea: 1979 Lamport one-time signatures.
- ▶ 1979 Merkle extends to more signatures.

Pros:

- Post quantum
- Only need secure hash function
- Security well understood
- Fast

Cons:

- Biggish signature though some tradeoffs possible
- Stateful, i.e., ever reusing a subkey breaks security.
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- Security well understood
- Fast
- We can count: OS update, code signing, ... naturally keep state.

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PROJECTS

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- NIST has gone through two rounds of requests for public input, most are positive and recommend standardizing XMSS and LMS. Only concern is about statefulness in general.

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Stateful Hash-Based Signatures

 ISO SC27 JTC1 WG2 has started a study period on stateful hash-based signatures.