Cryptography I, homework sheet 11 Due: 23 December 2011, 10:45

Both exercises can be done with the help of a computer but you should submit your programs as part of the homework solution. The program can be based on any computer algebra system, in particular for computing in \mathbb{F}_{1013}^* and \mathbb{F}_{1019}^* . Make sure that your programs compile and run correctly; my students will not debug your programs. The program should be humanly readable.

1. The schoolbook version of Pollard's rho method is often described with a simpler iteration function than we had in class.

Let $G_0 = g, b_0 = 1$, and $c_0 = 0$ and define

$$G_{i+1} = \begin{cases} G_i \cdot g \\ G_i^2 \\ G_i \cdot h \end{cases}, b_{i+1} = \begin{cases} b_i + 1 \\ 2b_i \\ b_i \end{cases}, c_{i+1} = \begin{cases} c_i \\ 2c_i \\ c_i + 1 \end{cases} \text{ for } G_i \equiv \begin{cases} 0 \mod 3 \\ 1 \mod 3 \\ 2 \mod 3 \end{cases}, a_i = \begin{cases} c_i \\ 1 \mod 3 \\ 2 \mod 3 \end{cases}$$

where one takes G_i as an integer.

Use this definition to attack the discrete logarithm problem given by g = 3, h = 245in \mathbb{F}_{1013}^* , i.e. find an integer 0 < a < 1012 such that $h = g^a$, using the G_i as defined above and $H_i = G_{2i}$.

Note that this version offers less randomness in the walk, splitting into more than 3 sets increases the randomness. The walk could start at any $G_0 = g^i h^j$ for known *i* and j – but then the homework would be harder to correct.

2. For a numerical example for the index calculus attack have a look at http://hyperelliptic.org/tanja/teaching/NTCrypto10/pictures/19-Nov-10, in particular IMGP2371.JPG.

Use factor base $\mathcal{F} = \{2, 3, 5, 7, 11, 13\}$ to solve the DLP h = 281, g = 2, in \mathbb{F}_{1019}^* .