# Cryptographic hash functions II Pollard rho for collision search

Tanja Lange

Eindhoven University of Technology

2MMC10 - Cryptology

## Generic hardness

The birthday paradox implies that if one draws elements at random from a set of *m* elements, then with 50% probability one has picked one element twice after about  $\sqrt{\pi m/2}$  picks. Hence it takes  $O(2^{n/2})$  calls to *H* to find a collision.

# Generic hardness

The birthday paradox implies that if one draws elements at random from a set of *m* elements, then with 50% probability one has picked one element twice after about  $\sqrt{\pi m/2}$  picks. Hence it takes  $O(2^{n/2})$  calls to *H* to find a collision.

Two problems need to be solved to use the rho method successfully:

- Design step function so that it "randomly" samples elements (so that the birthday paradox applies) while being deterministic (so we can use Floyd's cycle finding method to remove storage).
- 2 Design step function so that collision gives meaningful result.

## Generic hardness

The birthday paradox implies that if one draws elements at random from a set of *m* elements, then with 50% probability one has picked one element twice after about  $\sqrt{\pi m/2}$  picks. Hence it takes  $O(2^{n/2})$  calls to *H* to find a collision.

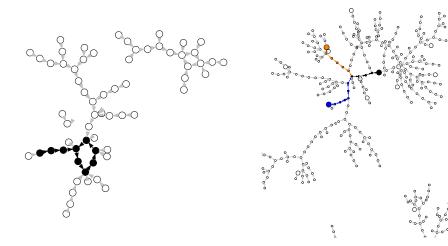
Two problems need to be solved to use the rho method successfully:

Design step function so that it "randomly" samples elements (so that the birthday paradox applies) while being deterministic (so we can use Floyd's cycle finding method to remove storage).

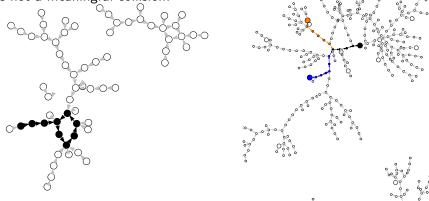
2 Design step function so that collision gives meaningful result. If message space includes output space, e.g.,  $H : \{0,1\}^* \to \{0,1\}^n$ , use random starting point  $W_0$  and iterate  $W_{i+1} = H(W_i)$ .

Else  $H : \mathbf{M} \to \mathbf{H}$  for some message space  $\mathbf{M}$  and hash space  $\mathbf{H}$ , compose with some map  $\phi : \mathbf{H} \to \mathbf{M}$ , iterate  $W_{i+1} = H(\phi(W_i))$ .

Can use parallel version due to van Oorshot and Wiener.



When the collision is detected (in Floyd or at distinguished point) it is not a meaningful collision!



When the collision is detected (in Floyd or at distinguished point) it is not a meaningful collision! Need to backtrack tilbcollision with different inputs

When the collision is detected (in Floyd or at distinguished point) it is not a meaningful collision!

Need to backtrack till collision with different inputs

For parallel the store length along with start and endpoint. Let W and W' reach the same point after  $\ell$  and  $\ell'$  steps ( $\ell < \ell'$ )

• Compute 
$$W'' = H^{\ell' - \ell}(W')$$
.

When the collision is detected (in Floyd or at distinguished point) it is not a meaningful collision!

Need to backtrack tilbcolfision with different inputs

For parallel the store length along with start and endpoint. Let W and W' reach the same point after  $\ell$  and  $\ell'$  steps ( $\ell < \ell'$ )

- Compute  $W'' = H^{\ell'-\ell}(W')$ . Unlucky if W'' = W.
- Iterate the following three steps until a collision is found:
  - Compare W and W''. If they are equal, output  $W_o$  and  $W''_o$ .
  - Store  $W_o = W, W''_o = W''$ .
  - Update W = H(W), W'' = H(W'').

Tanja Lange

Cryptographic hash functions II

## How to get interesting collisions?

We typically want meaningful collisions, e.g., colliding pdf files. This is one of the places where  $\mathbf{H} \not\subset \mathbf{M}$  happens.

Specify *n* bit positions in **M** that permit variations, make  $\phi$  map to those injectively. Several advanced formats are very flexible in taking macros or invisible images.

## How to get interesting collisions?

We typically want meaningful collisions, e.g., colliding pdf files. This is one of the places where  $\mathbf{H} \not\subset \mathbf{M}$  happens.

Specify *n* bit positions in **M** that permit variations, make  $\phi$  map to those injectively. Several advanced formats are very flexible in taking macros or invisible images.

Even more meaningful collisions

- Chosen prefix collision: want m = p\* and m' = p'\* to collide, where \* stands for arbitrary strings.
- Read Predicting the winner of the 2008 US Presidential
  Elections using a Sony PlayStation 3 for an impressive example of 12 pdf files all having the same hash, using chosen prefix collisions in MD5 and the flexibility of the pdf standard.